### Lightning watchtowers and why you should use them Sergi Delgado







# LIGHTNING 101





## CHANNEL LIFE CYCLE: OPENING A CHANNEL

- When two parties want to open a channel they lock-up funds in a 2-2 multisig contract (P2WSH)
- This lockup translates into a transaction on chain, known as funding transaction
- This is currently, mostly, single funded

from: A

- funding transaction
  - to: AB

## CHANNEL LIFE CYCLE: USING A CHANNEL I

- Each time sats are sent trough the channel a new transaction is created (actually two)
- All this transactions spend from the funding transaction and are known as commitment transactions
- Each side of the channel has a non-symmetrical version of each commitment
- Commitment transaction are time-locked as follows: - The spender has to wait to spend his funds (CSV) - The other side of the channel can spend straightaway

## CHANNEL LIFE CYCLE: USING A CHANNEL I

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	create		Commitment tra
-	All thi		$\sim$
	are kr	from: AB amount : x to_local amount: to_remo amount	to_local (A) <b>s</b> amount: y
-	Each		
	comm		to_remote (B) amount: x-y
-	Comr	APOV	
	- Th		
- The other side of the cha			



## CHANNEL LIFE CYCLE: USING A CHANNEL II

- A new commitment transaction revokes the previous one - All revoked commitments are valid Bitcoin transactions - Only a single commitment transaction is valid on chain - Only the owners of the channel know which is the last
- commitment

### commitment transactions



## CHANNEL LIFE CYCLE: CLOSING A CHANNEL

- The last commitment transaction is used as the base to create a
  - transaction without time locks
- This transaction is known as closing transaction
- Both parties can spend straightaway from the closing transaction
- The closing transaction will be stored on chain

from: AB

closing transaction

to: A to: B

## CHANNEL LIFE CYCLE: BREACHING A CHANNEL

- A side of the channel can try to spend an already revoked commitment (either "by mistake" or maliciously)
- This attempt is known as **channel breaching** and can result in the other side getting less funds than they deserve
- Only the channel owners can identify a revoked commitment, therefore **both sides may remain online**
- This is known as the always online assumption / requirement

## CHANNEL LIFE CYCLE: PUNISHING MISBEHAVIOUR I

- If a channel is breached by one side, the other side can (and will) penalise the attempt
- A transaction will be created using the revocation data and spending from the channel breach
- This transaction is known as **penalty transaction** and it claims all funds of the channel instantly



penalty transaction

## CHANNEL LIFE CYCLE: PUNISHING MISBEHAVIOUR II

# But what if the owners of the channel were not the only ones who could react to channel breaches?

## CHANNEL LIFE CYCLE: PUNISHING MISBEHAVIOUR II

# But what if the owners of the channel were not the only ones who could react to channel breaches?









## GENERAL CONCEPT

(**AKA Watchtowers**)?

User:

- Sends data to the server alongside a trigger condition

### Server:

- Looks for triggers on a communication channel

### What is the general paradigm behind third party watching systems

# - If the a trigger is seen, perform an action with the provided data





























[...] commitment\_txid, penalty\_tx, [...]





appointment

[...] commitment\_txid, penalty\_tx, [...]







appointment

[...] commitment\_txid, penalty\_tx, [...]







appointment

[...] commitment\_txid, penalty\_tx, [...]







appointment







appointment



## **BASIC WATCHTOWER PROTOCOL** appointment [...] commitment\_txid, penalty\_tx, commitment\_txid [...]









appointment







appointment







appointment



![](_page_29_Picture_1.jpeg)

![](_page_29_Picture_2.jpeg)

![](_page_29_Picture_4.jpeg)

![](_page_29_Figure_5.jpeg)

![](_page_30_Picture_1.jpeg)

![](_page_30_Picture_2.jpeg)

![](_page_30_Figure_4.jpeg)

![](_page_31_Picture_1.jpeg)

![](_page_31_Picture_2.jpeg)

![](_page_31_Picture_4.jpeg)

![](_page_31_Picture_5.jpeg)

![](_page_31_Figure_6.jpeg)

## WHY USING A TOWER (OR WHY NOT)? I

- High availability -
- Data redundancy for your node

of the properties it provides (e.g. mobile nodes)

So, when may we need one and what are the alternatives?

A watchtower is a failsafe mechanism with, mainly, two properties:

- It is useful for all kind of nodes, but specially for those that lack some

## WHY USING A TOWER (OR WHY NOT)? II

means:

- Power supply redundancy
- Internet service redundancy
- Data redundancy

We may **not need** a tower if we have a highly available node. That

### This may be the case of a top tier routing node, **not the average node.** Even here, some types of towers may be worth considering.

## WHY USING A TOWER (OR WHY NOT)? III

We may **need** a tower if we have:

- A non-highly available node (see previous slide) - A mobile phone node

average user is unlikely to run a highly available node.

Currently, the network is mostly run by techies and enthusiast, but the

![](_page_34_Picture_6.jpeg)

## WHY USING A TOWER (OR WHY NOT)? IV

Mobile nodes:

- Intermittent Internet access
- Lower bandwidth use w.r.t routing nodes
- Easier to lose data (phone breaks / gets stolen / ...)
- Node may not be always online
- Sporadically used (mainly when paying, can be offline for days)

## WHY USING A TOWER (OR WHY NOT)? IV

### Mobile nodes:

- Intermittent Int
- Lower bandwi
- Easier to lose
- Node may not
- Sporadically u

### **Background service**

It seems that Eclair Mobile has not been able to run in background lately. Make sure that your phone does not aggressively optimize this application.

Some vendors like **Nokia**, **Xiaomi**, **Huawei** or **Samsung** run overzealous custom battery savers preventing non white-listed apps to run in background.

You can check <u>our FAQ</u> for more information.

OK

### stolen / ...)

an be offline for days)

![](_page_37_Picture_0.jpeg)

![](_page_37_Picture_1.jpeg)

# Types of towers

![](_page_37_Picture_3.jpeg)

### For every channel update:

- The penalty transaction is **encrypted** under a key derived from the commitment transaction id
- A locator is also derived from the commitment transaction id
- The tower receives the encrypted blob and the locator

![](_page_38_Picture_5.jpeg)

Tadge Dryja - Unlinkable Outsourced Channel Monitoring - Scaling Bitcoin 2016

![](_page_38_Picture_8.jpeg)

<u>User side</u>

![](_page_39_Picture_3.jpeg)

<u>User side</u>

![](_page_40_Picture_4.jpeg)

### penalty\_tx = 020000000001010d8b7512b1f530338ca886...1f9624914fb8a680000000

User side

**commitment\_txid = 4a5e1e4baab89f3a32518...cc77ab2127b7afdeda33** 

![](_page_41_Picture_5.jpeg)

### penalty\_tx = 0200000000001010d8b7512b1f530338ca886...1f9624914fb8a680000000

User side

**commitment\_txid = 4a5e1e4baab89f3a32518...cc77ab2127b7afdeda33** 

**16 MSB** 

![](_page_42_Picture_6.jpeg)

### penalty\_tx = 0200000000001010d8b7512b1f530338ca886...1f9624914fb8a680000000

User side

**commitment\_txid = 4a5e1e4baab89f3a32518...cc77ab2127b7afdeda33** 

![](_page_43_Picture_6.jpeg)

### penalty\_tx = 0200000000001010d8b7512b1f530338ca886...1f9624914fb8a680000000

**16 MSB** → locator

### User side

**commitment\_txid = 4a5e1e4baab89f3a32518...cc77ab2127b7afdeda33** 

### cipher = CHACHA20POLY1305 sk = SHA256(commitment\_txid) IV = 0

![](_page_44_Picture_6.jpeg)

### penalty\_tx = 0200000000001010d8b7512b1f530338ca886...1f9624914fb8a680000000

16 MSB  $\longrightarrow$  locator

### User side

penalty\_tx = 020000000001010d8b7512b1f530338ca886...1f9624914fb8a680000000

**commitment\_txid = 4a5e1e4baab89f3a32518...cc77ab2127b7afdeda33** 16 MSB  $\longrightarrow$  locator

cipher = CHACHA20POLY1305 encrypt (penalty\_tx, sk, IV) sk = SHA256(commitment\_txid) IV = 0

![](_page_45_Picture_6.jpeg)

### User side

penalty\_tx = 0200000000001010d8b7512b1f530338ca886...1f9624914fb8a6800000000

**commitment\_txid = 4a5e1e4baab89f3a32518...cc77ab2127b7afdeda33** 16 MSB  $\longrightarrow$  locator

cipher = CHACHA20POLY1305 encrypt (penalty\_tx, sk, IV) sk = SHA256(commitment\_txid) IV = 0

![](_page_46_Picture_6.jpeg)

![](_page_46_Picture_7.jpeg)

### User side

**commitment\_txid = 4a5e1e4baab89f3a32518...cc77ab2127b7afdeda33 16 MSB** locator

cipher = CHACHA20POLY1305 encrypt (penalty\_tx, sk, IV) sk = SHA256(commitment\_txid) IV = 0

![](_page_47_Picture_6.jpeg)

### penalty\_tx = 020000000001010d8b7512b1f530338ca886...1f9624914fb8a680000000

![](_page_47_Picture_8.jpeg)

![](_page_47_Picture_9.jpeg)

### User side

penalty\_tx = 0200000000001010d8b7512b1f530338ca886...1f9624914fb8a6800000000

**commitment\_txid = 4a5e1e4baab89f3a32518...cc77ab2127b7afdeda33 16 MSB** locator

cipher = CHACHA20POLY1305 encrypt (penalty\_tx, sk, IV) sk = SHA256(commitment\_txid) IV = 0

![](_page_48_Picture_6.jpeg)

![](_page_48_Picture_7.jpeg)

### **SEND TO THE TOWER**

![](_page_48_Picture_9.jpeg)

Tower side

![](_page_49_Picture_2.jpeg)

Tower side

for every transaction\_id in every block

**locator = 16 MSB transaction\_id** 

![](_page_50_Picture_4.jpeg)

- Tower side
- for every transaction\_id in every block
  - **locator = 16 MSB transaction id** 
    - if locator in appointments:
      - sk = SHA256(transaction\_id)
      - IV = 0

![](_page_51_Picture_8.jpeg)

![](_page_51_Picture_9.jpeg)

- Tower side
- for every transaction\_id in every block
  - **locator = 16 MSB transaction\_id** 
    - if locator in appointments:
      - sk = SHA256(transaction\_id)
      - IV = 0

decrypt (encrypted blob, sk, IV)

![](_page_52_Picture_9.jpeg)

- Tower side
- for every transaction\_id in every block
  - **locator = 16 MSB transaction id** 
    - if **locator** in appointments:
      - sk = SHA256(transaction\_id)
      - IV = 0

decrypt (encrypted blob, sk, IV)

![](_page_53_Picture_9.jpeg)

![](_page_53_Picture_11.jpeg)

### **Pros:**

- Privacy preserving
- towers)
- No infrastructure needed for the user

### **Cons:**

- Design is rather complex
- O(N) storage
- Can be easily spammed (altruistic vs non-altruistic)

### Can give service to the whole network (with one or multiple

## PERSONAL USE ONLY TOWER I

If the tower is for personal use only, the design can be highly simplified:

- Privacy is not a concern (at least not at the same level)
- Data may not need to be encrypted -
- Most of the spam protections included in the design can be lifted
- Storage can be O(1)

![](_page_55_Picture_7.jpeg)

## PERSONAL USE ONLY TOWER II

**Examples:** 

- Channel ids can be shared with the tower, so no more need for locators (only renovation data) (**O(N) but simpler design**) - Even revocation keys could be shared, so the storage is drastically reduced (**O(1) but riskier**)

proposing / pursuing them though (h/t Antoine Riard)

Haven't seen any tower like this in the wild. I've seen some devs

## PERKS OF USING ANY KIND OF TOWER I

### CSV in to\_local outputs COULD be reduced:

CS

### c-lightning

### LND

Eclair

rust-lightning

SV (default)	~ time	
144	1D	
44-2016	1D - 2W	
720	5D	
144	1D	

## PERKS OF USING ANY KIND OF TOWER II

![](_page_58_Picture_1.jpeg)

Unresponsive peers. If you have to close unilaterally you are forced to wait the CSV delay.

However

It will be unsafe to set a too low CSV if no-one is watching your channels

![](_page_58_Figure_5.jpeg)

## PERKS OF USING ANY KIND OF TOWER III

![](_page_59_Picture_1.jpeg)

Having a too big CSV may hurt UX

Finally

the tower to respond

### This is not tradeoff free. Depending on the setup you may be trusting

![](_page_59_Picture_6.jpeg)

![](_page_60_Picture_0.jpeg)

![](_page_60_Picture_2.jpeg)

![](_page_60_Picture_3.jpeg)

## SPECIAL THANKS TO MY SPONSORS

- Square Crypto

8 sponsors are funding sr-gi's work.

![](_page_60_Picture_8.jpeg)

![](_page_61_Picture_0.jpeg)

![](_page_61_Picture_1.jpeg)

![](_page_61_Picture_2.jpeg)